



# Improvements in NTRU+

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## ❖ NTRU-based IND-CCA secure KEM &amp; PKE

GenNTRU[ $\psi_1^n$ ]OW-CPA secure  
PKEnegligible  
average-case  
correctness errorACWC<sub>2</sub> $\approx$ IND-CPA secure  
PKEnegligible  
worst-case  
correctness errorFO<sub>KEM</sub>IND-CCA secure  
KEMIND-CCA secure KEM  
with Re-encryption $\overline{\text{FO}}_{\text{KEM}}$ 

NTRU+KEM

IND-CCA secure KEM  
without Re-encryptionFO<sub>PKE</sub>IND-CCA secure  
PKEIND-CCA secure PKE  
with Re-encryption $\overline{\text{FO}}_{\text{PKE}}$ 

NTRU+PKE

IND-CCA secure PKE  
without Re-encryption

## ❖ Goal

- ◆ Improve the efficiency of the key generation algorithm !

Algorithm	sec. (c)	n	q	PK (Byte)	CT (Byte)	SK (Byte)	$\log_2 \delta$	Optimized C (K Cycles)			AVX2 (K Cycles)		
								Gen	Encap	Decap	Gen	Encap	Decap
NTRU+KEM 576	114	576	3,457	864	864	1,760	-487	274 ↓ 167	102 ↓ 80	128 ↓ 95	21 ↓ 24	23 ↓ 22	14 ↓ 13
NTRU+KEM 768	164	768		1,152	1,152	2,336	-379	325 ↓ 192	133 ↓ 101	172 ↓ 121	26 ↓ 26	30 ↓ 27	19 ↓ 16
NTRU+KEM 864	189	864		1,296	1,296	2,624	-340	314 ↓ 238	157 ↓ 123	207 ↓ 148	24 ↓ 28	32 ↓ 30	21 ↓ 19
NTRU+KEM 1152	263	1,152		1,728	1,728	3,488	-260	751 ↓ 370	198 ↓ 162	274 ↓ 196	45 ↓ 41	39 ↓ 39	25 ↓ 26

## ❖ Changes in the Specification & Implementation

### ◆ 1. Hash function

- Replaced AES256CTR with SHAKE256 for hash function instantiation

### ◆ 2. Key generation algorithm

- Samples the polynomials  $f$  and  $g$  separately
- Uses early abort in the invertibility test

### ◆ 3. Ring operation

- Modified the NTT structures for  $\text{NTRU}+\{\text{KEM}, \text{PKE}\}_{576, 768, 1152}$
- Applied the Montgomery reduction lazily
- Improved the modulus inversion

❖ Samples the polynomials  $f$  and  $g$  separately**Algorithm 15**  $\text{Gen}(1^\lambda)$ : key generation**Ensure:** Public key  $pk \in \mathcal{B}^{\lceil \log_2 q \rceil \cdot n/8}$ **Ensure:** Secret key  $sk \in \mathcal{B}^{\lceil \log_2 q \rceil \cdot n/4}$ 

```

1:  $d \leftarrow \mathcal{B}^{32}$ 
2:  $(f, g) := \text{XOF}(d, n/2)$ 
3:  $\mathbf{f}' := \text{CBD}_1(f)$ ,  $\mathbf{g}' := \text{CBD}_1(g)$ 
4:  $\mathbf{f} = 3\mathbf{f}' + 1$ 
5:  $\mathbf{g} = 3\mathbf{g}'$ 
6:  $\hat{\mathbf{f}} = \text{NTT}(\mathbf{f})$ ,  $\hat{\mathbf{g}} = \text{NTT}(\mathbf{g})$ 
7: if  $\mathbf{f}$  or  $\mathbf{g}$  is not invertible in  $R_q$ , restart
8:  $\hat{\mathbf{h}} = \hat{\mathbf{g}} \circ \hat{\mathbf{f}}^{-1}$ 
9:  $pk := \text{Encode}_q(2^{16} \cdot \hat{\mathbf{h}})$ 
10:  $sk := \text{Encode}_q(\hat{\mathbf{f}}) || \text{Encode}_q(2^{16} \cdot \hat{\mathbf{h}}^{-1}) || F(pk)$ 
11: return  $(pk, sk)$ 

```

**Algorithm 15:**  $\text{Gen}(1^\lambda)$ : key generation**Ensure:** Public key  $pk \in \mathcal{B}^{\lceil \log_2 q \rceil \cdot n/8}$ **Ensure:** Secret key  $sk \in \mathcal{B}^{\lceil \log_2 q \rceil \cdot n/4}$ 

```

1: repeat
2:    $d \leftarrow \mathcal{B}^{32}$ 
3:    $f := \text{XOF}(d, n/4)$ 
4:    $\mathbf{f}' := \text{CBD}_1(f)$ 
5:    $\mathbf{f} := 3\mathbf{f}' + 1$ 
6:    $\hat{\mathbf{f}} := \text{NTT}(\mathbf{f})$ 
7: until  $\mathbf{f}$  is invertible in  $R_q$ 
8: repeat
9:    $d \leftarrow \mathcal{B}^{32}$ 
10:   $g := \text{XOF}(d, n/4)$ 
11:   $\mathbf{g}' := \text{CBD}_1(g)$ 
12:   $\mathbf{g} := 3\mathbf{g}'$ 
13:   $\hat{\mathbf{g}} := \text{NTT}(\mathbf{g})$ 
14: until  $\mathbf{g}$  is invertible in  $R_q$ 
15:  $\hat{\mathbf{h}} := \hat{\mathbf{g}} \circ \hat{\mathbf{f}}^{-1}$ 
16:  $pk := \text{Encode}_q(2^{16} \cdot \hat{\mathbf{h}})$ 
17:  $sk := \text{Encode}_q(\hat{\mathbf{f}}) || \text{Encode}_q(2^{16} \cdot \hat{\mathbf{h}}^{-1}) || F(pk)$ 
18: return  $(pk, sk)$ 

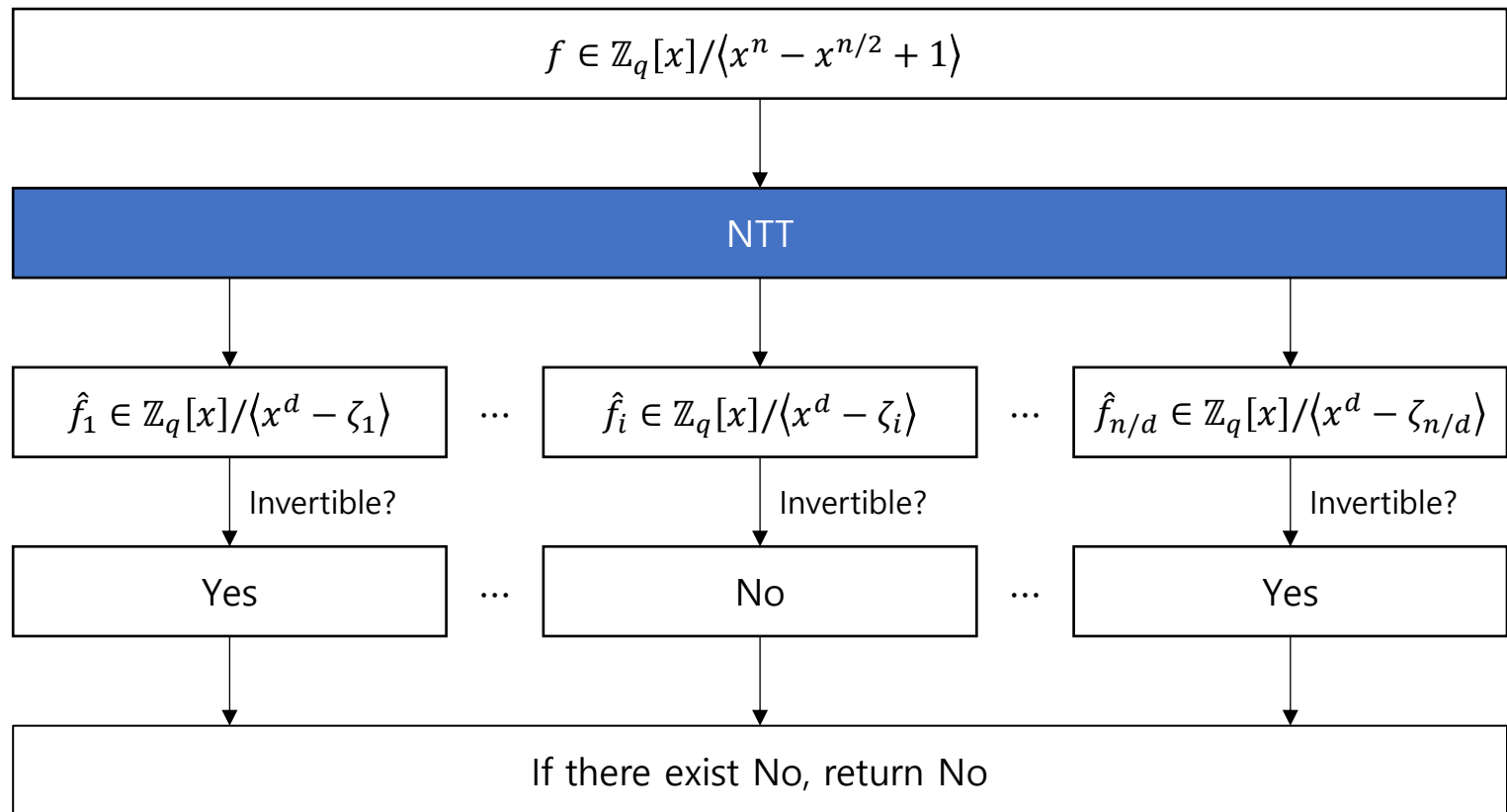
```

Version 2.0

Version 2.2

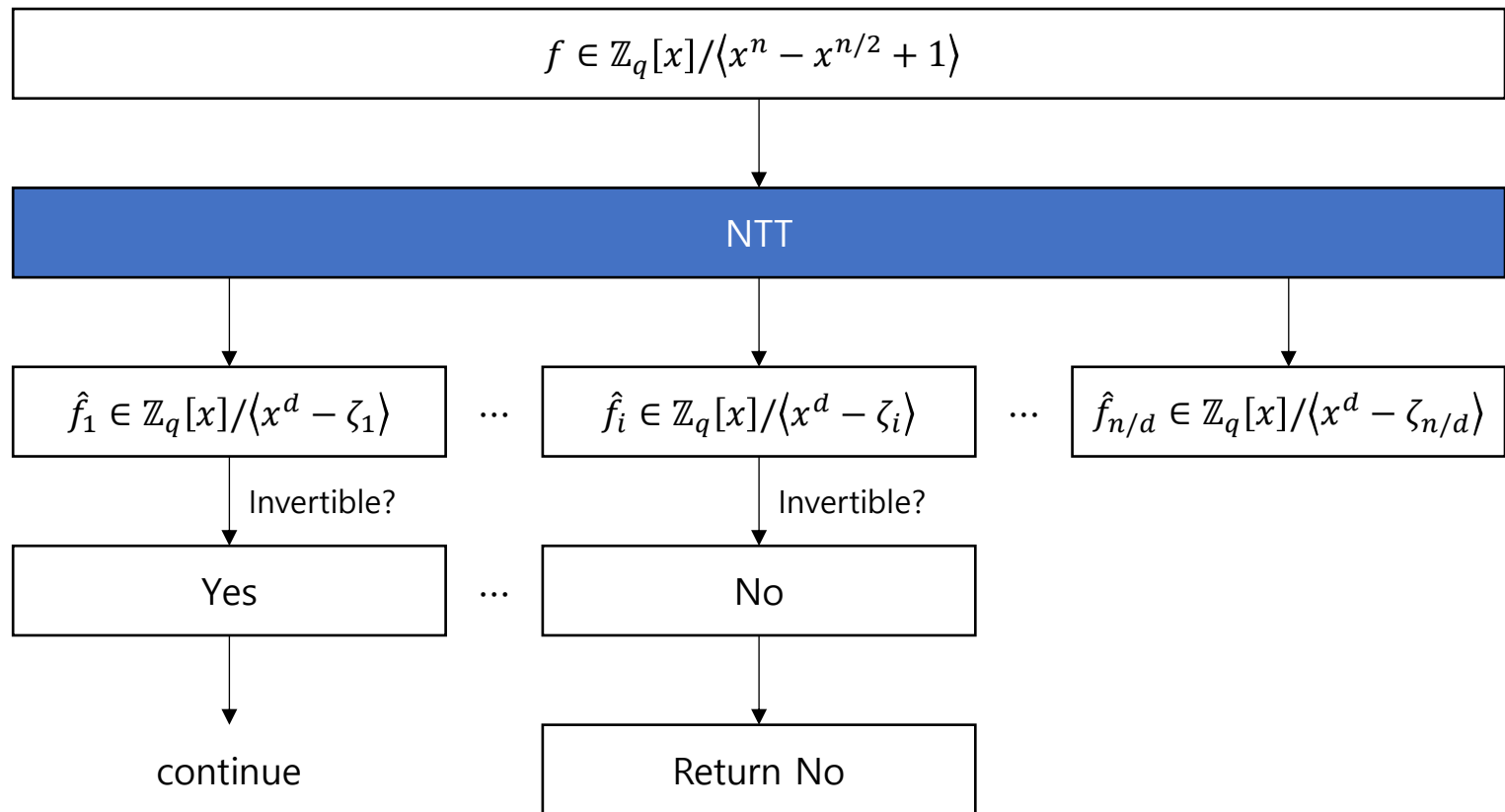
## ❖ Uses early abort in the invertibility test

## ◆ Invertibility Test



## ❖ Uses early abort in the invertibility test

## ◆ Invertibility Test with Early Abort



## ❖ Modified the NTT structures for NTRU+{KEM, PKE}{576, 768, 1152}

- ◆ NTT structure for  $\mathbb{Z}_{3457}[x]/\langle x^n - x^{n/2} + 1 \rangle$  with  $n = 576$ 
  - Similar analysis can be applied to  $n = 768, 1152$

	Naïve	Version 1.0	Version 1.1	Version 2.2
Radix-2 NTT layers	6	5	6	4
Radix-3 NTT layers	2	2	1	2
Result of NTT	$\prod_{i=1}^{576} \mathbb{Z}_{3457}[x]/\langle x - \psi_i \rangle$	$\prod_{i=1}^{288} \mathbb{Z}_{3457}[x]/\langle x^2 - \eta_i \rangle$	$\prod_{i=1}^{192} \mathbb{Z}_{3457}[x]/\langle x^3 - \xi_i \rangle$	$\prod_{i=1}^{144} \mathbb{Z}_{3457}[x]/\langle x^4 - \xi_i \rangle$
Polynomial inversion	576 modulus inversions	288 modulus inversions	192 modulus inversions	144 modulus inversions
Precomputation table	576 elements	288 elements	192 elements	144 elements

- The ideas from [ZFY23] were adapted to implement efficient inversion in  $\mathbb{Z}_{3457}[x]/\langle x^4 - \xi_i \rangle$ .
  - It computes the inversion in  $\mathbb{Z}_{3457}[x]/\langle x^n - \xi_i \rangle$  by using the inversion in  $\mathbb{Z}_{3457}[x]/\langle x^{n/2} - \xi_i \rangle$



### ❖ Applied the Montgomery reduction lazily

#### ◆ Multiplication in $R_q = \mathbb{Z}_q[x]/\langle x^3 - \zeta \rangle$

- $a(x) = a_0 + a_1x + a_2x^2 \in R_q$
- $b(x) = b_0 + b_1x + b_2x^2 \in R_q$
- $c(x) = a(x)b(x)$ 
  - $c_0 = a_0b_0 + (a_1b_2 + a_2b_1)\zeta$
  - $c_1 = a_0b_1 + a_1b_0 + a_2b_2\zeta$
  - $c_2 = a_0b_2 + a_1b_1 + a_2b_0$

#### ◆ Multiplication with Montgomery & Barrett reduction

$$a_0b_2 + a_1b_1 + a_2b_0$$



$$\text{barrett}(\text{mont}(a_0 \cdot b_2) + \text{mont}(a_1 \cdot b_1) + \text{mont}(a_2 \cdot b_0))$$



$$\text{mont}(a_0 \cdot b_2 + a_1 \cdot b_1 + a_2 \cdot b_0)$$

Apply Montgomery reduction lazily

**This technique is already applied in NCC-Sign Version 2.1**

## ❖ Improved the modulus inversion

## ◆ Modulus inversion

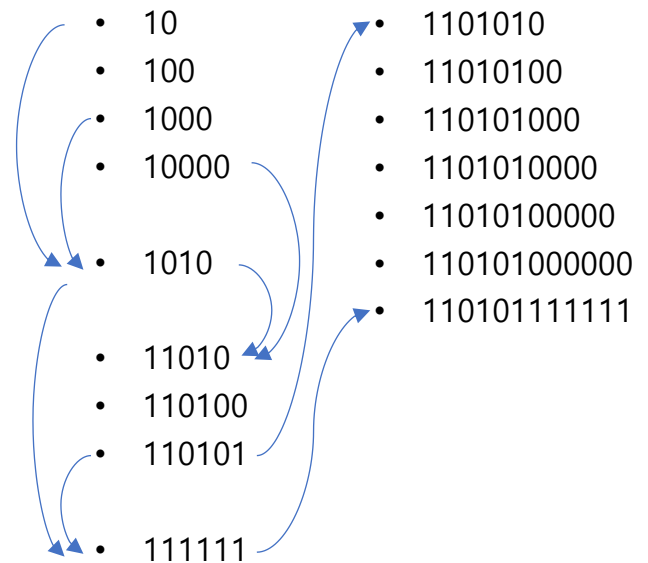
$$a^{-1} \equiv a^{q-2} \pmod{q}, \quad q - 2 = 3455 = 110101111111_{(2)}$$

## ▪ Modulus inversion with Square-and-Multiply

- |           |                |
|-----------|----------------|
| • 10      | • 11010110     |
| • 11      | • 11010111     |
| • 110     | • 110101110    |
| • 1100    | • 110101111    |
| • 1101    | • 1101011110   |
| • 11010   | • 1101011111   |
| • 110100  | • 11010111110  |
| • 110101  | • 11010111111  |
| • 1101010 | • 110101111110 |
| • 1101011 | • 110101111111 |

20 multiplications

## ▪ Improved Modulus inversion



16 multiplications

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Kyber512	117	256x2	3,329	800	768	1,632	-139	116	137	158	36	39	24
Kyber768	181	256x3		1,184	1,088	2,400	-164	182	202	230	51	55	37
Kyber1024	253	256x4		1,568	1,568	3,168	-174	270	321	359	65	73	52
NTRUHPS 2048509	106	509	2,048	699	699	935	$-\infty$	8,031	746	1,384	376	262	33
NTRUHRSS 701	136	701	8,192	1,138	1,138	1,450	$-\infty$	14,684	1030	2,623	365	166	52
NTRUHPS 2048677	145	677	2,048	930	930	1,234	$-\infty$	13,882	1206	2,441	545	348	49
NTRUHPS 4096821	179	821	4096	1,230	1,230	1,590	$-\infty$	20,385	1,644	3,519	702	423	62

**T**hank You

**Q&A**